

矩形网格上的二元切触有理插值*

荆科^{1,2}, 康宁²

- (1. 阜阳师范学院 数学与统计学院, 安徽 阜阳 236037;
2. 合肥工业大学 管理学院, 合肥 230009)

摘要: 二元切触有理插值是有理插值的一个重要内容,而降低其函数的次数和解决其函数的存在性是有理插值的一个重要问题.二元切触有理插值算法的可行性大都是有条件的,且计算复杂度较大,有理函数的次数较高.利用二元 Hermite(埃米特)插值基函数的方法和二元多项式插值误差性质,构造出了一种二元切触有理插值算法并将其推广到向量值情形.较之其它算法,有理插值函数的次数和计算量较低.最后通过数值实例说明该算法的可行性是无条件的,且计算量低.

关键词: 二元切触有理插值; 误差公式; 二元 Hermite 插值; 基函数

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引 言

切触有理插值是类似于多项式 Hermite 插值的一种插值.二元切触有理插值问题是一元切触有理插值问题的自然推广.由于其插值节点的复杂性,它比一元情形要困难得多.目前关于切触有理插值的研究主要集中在一元情形.文献[1]将切触有理插值的非线性问题转化为线性问题.文献[2]证明了切触有理插值函数的存在性及唯一性.文献[3]利用连分式的方法给出了一种经典的切触有理插值算法.文献[4-5]利用 Newton-Padé 逼近方法给出了一种切触有理插值算法.文献[6]利用 Hermite-Newton 插值公式给出了一个判断切触有理插值函数存在性的充要条件.文献[7]给出了切触有理插值的 Neville(内维尔)递推算法.文献[8]利用 Hermite 基函数和多项式插值误差公式的方法,给出了一种切触有理插值算法,降低了有理插值函数的次数.关于二元切触有理插值的研究较少,且大都针对矩形网格上的切触有理函数的构造方法.所谓的二元切触有理插值问题就是在矩形网格

$$\Pi^{m,n} = \{ (x_i, y_j) / (x_i, y_j) \in R^2, i = 0, 1, \dots, s; j = 0, 1, \dots, t \} \quad (1)$$

上寻求二元有理函数 $r(x, y) = p(x, y)/q(x, y)$, 使之满足下列插值条件:

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作者简介: 荆科(1983—),男,安徽颖上人,博士生(E-mail: jingxuefei296@sina.com);

康宁(1986—),女,安徽利辛人,讲师,博士生(通讯作者. E-mail: kangningning2006@126.com).

$$\left. \frac{\partial^{u+v}}{\partial x^u \partial y^v} \left(\frac{p(x,y)}{q(x,y)} \right) \right|_{(x_i, y_j)} = f_{ij}^{(u,v)}, \quad (2)$$

其中 $u = 0, 1, \dots, m_i - 1; v = 0, 1, \dots, n_j - 1$. $p(x, y), q(x, y)$ 均是二元多项式.

文献[9]利用文献[10-11]中 Newton 插值的思想和分段组合的方法,给出了一种矩形网格上的二元切触有理插值算法,不仅算法的可行性是无条件的,且有理函数次数较低.文献[12]给出了一种混合 Thiele-Werner 型切触有理插值算法,但有理函数次数较高.文献[13]利用 Salzer 型插值连分式和扩展的 Newton 插值多项式构造的算法,具有很好的逼近效果.文献[14]将在切触有理插值中起重要作用的 Salzer 定理推广到了多元向量的情形.上述结果虽然很好,但算法的可行性大都是有条件的,且计算复杂度较高.本文利用二元多项式插值的误差性质及 Hermite 基函数的方法,建立一种二元有理函数插值公式,并将其推广到向量值情形.较之其它算法,不仅算法的可行性是无条件的,且具有有理函数次数较低、计算量较小等特点.

1 切触有理插值公式

下面就 $m_i = 2, i = 0, 1, \dots, m; n_j = 2, j = 0, 1, \dots, n$ 的情形进行讨论.

设 $\alpha_i(x), \beta_i(x)$ 分别表示一元 Hermite 插值对于变量 x 的基函数,同时 $\alpha_j(y), \beta_j(y)$ 表示对于变量 y 的基函数,且满足下列条件:

$$\begin{cases} \alpha_i(x_k) = \delta_{ik} = \begin{cases} 0, & i \neq k, \\ 1, & i = k, \end{cases} & \alpha'_i(x_k) = 0, \\ \beta_i(x_k) = 0, & \beta'_i(x_k) = \delta_{ik}, \quad i, k = 0, 1, \dots, s, \end{cases} \quad (3)$$

$$\begin{cases} \alpha_j(y_l) = \delta_{jl} = \begin{cases} 0, & j \neq l, \\ 1, & j = l, \end{cases} & \alpha'_j(y_l) = 0, \\ \beta_j(y_l) = 0, & \beta'_j(y_l) = \delta_{jl}, \quad j, l = 0, 1, \dots, t. \end{cases} \quad (4)$$

定义有理函数如下:

$$\lambda_{ij}(x, y) = \frac{q_{ij} \alpha_i(x) \alpha_j(y) + q_{ij}^{(1,0)} \alpha_j(y) \beta_i(x) + q_{ij}^{(0,1)} \alpha_i(x) \beta_j(y) + q_{ij}^{(1,1)} \beta_i(x) \beta_j(y)}{q(x, y)},$$

其中

$$q(x, y) = \sum_{i=0}^s \sum_{j=0}^t [q_{ij} \alpha_i(x) \alpha_j(y) + q_{ij}^{(1,0)} \alpha_j(y) \beta_i(x) + q_{ij}^{(0,1)} \alpha_i(x) \beta_j(y) + q_{ij}^{(1,1)} \beta_i(x) \beta_j(y)].$$

直接验证 $\lambda_{ij}(x, y)$ 具有如下性质:

$$1) \lambda_{ij}(x, y) \text{ 是 } [(2s + 1, 2t + 1)/(2s + 1, 2t + 1)] \text{ 型有理函数}; \quad (5)$$

$$2) \lambda_{ij}(x_k, y_l) = \begin{cases} 1, & i = k, \text{ and } j = l \\ 0, & i \neq k, \text{ or } j \neq l \end{cases} (i, k = 0, 1, \dots, s; j, l = 0, 1, \dots, t); \quad (6)$$

$$3) \frac{\partial \lambda_{ij}(x_k, y_l)}{\partial x} = \frac{\partial \lambda_{ij}(x_k, y_l)}{\partial y} = \frac{\partial^2 \lambda_{ij}(x_k, y_l)}{\partial x \partial y} = 0 \quad (i, k = 0, 1, \dots, s; j, l = 0, 1, \dots, t). \quad (7)$$

先讨论二元数量值切触有理插值问题.引入插值算子如下:

$$p_{ij}(x, y) = (x - x_i)(y - y_j)f_{ij}^{(1,1)} + (x - x_i)f_{ij}^{(1,0)} + (y - y_j)f_{ij}^{(0,1)} + f_{ij}. \quad (8)$$

显然

$$p_{ij}(x_i, y_j) = f_{ij}, \quad \frac{\partial p_{ij}(x_i, y_j)}{\partial x} = f_{ij}^{(1,0)}, \quad \frac{\partial p_{ij}(x_i, y_j)}{\partial y} = f_{ij}^{(0,1)}, \quad \frac{\partial^2 p_{ij}(x_i, y_j)}{\partial x \partial y} = f_{ij}^{(1,1)}$$

$$(i = 0, 1, \dots, s; j = 0, 1, \dots, t).$$

定理 1 对于给定插值节点(1)和相应函数值 f_{ij} 及其偏导数值 $f_{ij}^{(u,v)}$,

$$r(x, y) = \sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x, y) p_{ij}(x, y) \quad (9)$$

是满足插值条件的二元切触有理插值函数.

证明

$$r(x_k, y_l) = \sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x_k, y_l) p_{ij}(x_k, y_l) = p_{kl}(x_k, y_l) = f_{kl},$$

$$\frac{\partial r(x_k, y_l)}{\partial x} = \sum_{i=0}^s \sum_{j=0}^t \left[\frac{\partial \lambda_{ij}(x_k, y_l)}{\partial x} p_{ij}(x_k, y_l) + \lambda_{ij}(x_k, y_l) \frac{\partial p_{ij}(x_k, y_l)}{\partial x} \right] =$$

$$\sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x_k, y_l) \frac{\partial p_{ij}(x_k, y_l)}{\partial x} = \frac{\partial p_{kl}(x_k, y_l)}{\partial x} = f_{kl}^{(1,0)},$$

$$\frac{\partial r(x_k, y_l)}{\partial y} = \sum_{i=0}^s \sum_{j=0}^t \left[\frac{\partial \lambda_{ij}(x_k, y_l)}{\partial y} p_{ij}(x_k, y_l) + \lambda_{ij}(x_k, y_l) \frac{\partial p_{ij}(x_k, y_l)}{\partial y} \right] =$$

$$\sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x_k, y_l) \frac{\partial p_{ij}(x_k, y_l)}{\partial y} = \frac{\partial p_{kl}(x_k, y_l)}{\partial y} = f_{kl}^{(0,1)},$$

$$\frac{\partial^2 r(x_k, y_l)}{\partial x \partial y} = \sum_{i=0}^s \sum_{j=0}^t \left[\frac{\partial^2 \lambda_{ij}(x_k, y_l)}{\partial x \partial y} p_{ij}(x_k, y_l) + \lambda_{ij}(x_k, y_l) \frac{\partial^2 p_{ij}(x_k, y_l)}{\partial x \partial y} + \right.$$

$$\left. \frac{\partial \lambda_{ij}(x_k, y_l)}{\partial x} \frac{\partial p_{ij}(x_k, y_l)}{\partial y} + \frac{\partial \lambda_{ij}(x_k, y_l)}{\partial y} \frac{\partial p_{ij}(x_k, y_l)}{\partial x} \right] =$$

$$\sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x_k, y_l) \frac{\partial^2 p_{ij}(x_k, y_l)}{\partial x \partial y} = \frac{\partial^2 p_{kl}(x_k, y_l)}{\partial x \partial y} = f_{kl}^{(1,1)}$$

$$(i = 0, 1, \dots, s; j = 0, 1, \dots, t).$$

可以看出由式(9)得到的二元切触有理函数满足插值条件,但缺点是有理函数的次数太高.为了解决这个问题,引入下列记号:

$$f_{vu} = f(x_0, x_1, \dots, x_v; y_0, y_1, \dots, y_u).$$

当 $v = 2k, k \in \mathbf{Z}$

$$f_{vu} = f(x_0, x_0, x_1, x_1, \dots, x_k, x_k; y_0, y_1, \dots, y_u),$$

$$x_v = (x - x_0)^2(x - x_1)^2 \cdots (x - x_k)^2.$$

当 $v = 2k + 1, k \in \mathbf{Z}$

$$f_{vu} = f(x_0, x_0, x_1, x_1, \dots, x_k, x_k, x_{k+1}; y_0, y_1, \dots, y_u),$$

$$x_v = (x - x_0)^2(x - x_1)^2 \cdots (x - x_k)^2(x - x_{k+1}).$$

同样对于 $u = 2k, k \in \mathbf{Z}, u = 2k + 1, k \in \mathbf{Z}$ 也是一样.

用文献[15]中的二元多项式插值公式可类似证明下述定理.

定理 2 矩形网格(1)上满足函数值 f_{ij} 及其偏导数值 $f_{ij}^{(u,v)}$ 的二元 Hermite 多项式插值公式为

$$f(x, y) = \sum_{v=0}^{2s} \sum_{u=0}^{2t} f_{vu} x_v y_u + r(x, y), \quad (10)$$

其中

$$e(x, y) = \frac{\omega_{2s+1}(x)}{(2s+1)!} f_x^{(2s+1)}(\xi, y) + \frac{\omega_{2t+1}(y)}{(2t+1)!} f_x^{(2t+1)}(x, \eta) - \frac{\omega_{2s+1}(x)\omega_{2t+1}(y)}{(2s+1)!(2t+1)!} f_{xy}^{(2s+1, 2t+1)}(\bar{\xi}, \bar{\eta}). \quad (11)$$

此处 ξ 和 $\bar{\xi}$ 落在包含 x, x_0, x_1, \dots, x_s 的最小区间内,而 η 和 $\bar{\eta}$ 落在包含 y, y_0, y_1, \dots, y_t 的最小区间内.

证明 根据一元 Newton 差商插值公式有

$$f(x, y) = \sum_{v=0}^{2s} x_v f(x_0, x_1, \dots, x_v; y) + x_{2s+1} f(x, x_0, x_1, \dots, x_v; y),$$

$$f(x_0, x_1, \dots, x_v; y) = \sum_{u=0}^{2t} f_{vu} y_u + y_{2t+1} f(x_0, x_1, \dots, x_v; y, y_0, y_1, \dots, y_{2t}).$$

则

$$f(x, y) = \sum_{v=0}^{2s} \sum_{u=0}^{2t} f_{vu} x_v y_u + r(x, y),$$

其中

$$e(x, y) = x_{2s+1} f(x, x_0, x_1, \dots, x_{2s}; y) + y_{2t+1} f(x; y, y_0, y_1, \dots, y_{2t}) - x_{2s+1} y_{2t+1} f(x, x_0, x_1, \dots, x_{2s}; y, y_0, y_1, \dots, y_{2t}) - x_{2s+1} y_{2t+1} f(x, x_0, x_1, \dots, x_{2s}; y, y_0, y_1, \dots, y_{2t}).$$

推论 1 如果 $f(x, y)$ 是 x, y 的次数分别不超过 $2s+1, 2t+1$ 的二元代数多项式,则它在矩形网格(1)上的二元 Hermite 插值多项式 $p(x, y) \in P_{2s+1, 2t+1}$ 恒等于 $f(x, y)$;有时也说插值是精确的.

根据文献[16]的推论 1,类似地可以证明上述推论 1.

为了降低切触有理插值函数的次数和解决切触有理插值函数的存在性问题,本文设

$$q(x, y) = \sum_{i=0}^m \sum_{j=0}^n \omega_i(x) \omega_j(y),$$

因为当 $k=m, \dots, s; l=n, \dots, t, \omega_{ij}(x_k, y_l) > 0$,否则 $\omega_{ij}(x_k, y_l) = 1$,所以 $q(x, y)$ 在矩形网格(1)上的函数值 q_{ij} 大于 0.设 $q_{ij}, q_{ij}^{(1,0)}, q_{ij}^{(0,1)}, q_{ij}^{(1,1)}$ 是二元多项式 $q(x, y)$ 在矩形网格(1)上的函数值和偏导数值,根据推论 1, $q(x, y)$ 在矩形网格(1)上的二元 Hermite 插值多项式就是本身,其中 $m \leq 2s+1, n \leq 2t+1$ 为正整数,所以可得 $r(x, y)$ 是 $[(2s+2, 2t+2)/(m, n)]$ 型有理函数.且

$$r(x, y) = \sum_{i=0}^s \sum_{j=0}^t [q_{ij} \alpha_i(x) \alpha_j(y) + q_{ij}^{(1,0)} \alpha_j(y) \beta_i(x) + q_{ij}^{(0,1)} \alpha_i(x) \beta_j(y) + q_{ij}^{(1,1)} \beta_i(x) \beta_j(y)] p_{ij}(x, y) / \sum_{i=0}^s \sum_{j=0}^t [q_{ij} \alpha_i(x) \alpha_j(y) +$$

$$\begin{aligned}
 & [q_{ij}^{(1,0)}\alpha_j(y)\beta_i(x) + q_{ij}^{(0,1)}\alpha_i(x)\beta_j(y) + q_{ij}^{(1,1)}\beta_i(x)\beta_j(y)] = \\
 & \sum_{i=0}^s \sum_{j=0}^t [q_{ij}\alpha_i(x)\alpha_j(y) + q_{ij}^{(1,0)}\alpha_j(y)\beta_i(x) + q_{ij}^{(0,1)}\alpha_i(x)\beta_j(y) + \\
 & q_{ij}^{(1,1)}\beta_i(x)\beta_j(y)]p_{ij}(x,y) \Big/ \sum_{i=0}^m \sum_{j=0}^n \omega_i(x)\omega_j(y) = \\
 & \sum_{i=0}^s \sum_{j=0}^t [q_{ij}\alpha_i(x)\alpha_j(y) + q_{ij}^{(1,0)}\alpha_j(y)\beta_i(x) + q_{ij}^{(0,1)}\alpha_i(x)\beta_j(y) + \\
 & q_{ij}^{(1,1)}\beta_i(x)\beta_j(y)]p_{ij}(x,y) \Big/ [a_{m,n}x^m y^n + a_{m,n-1}x^m y^{n-1} + a_{m-1,n}x^{m-1} y^n + \\
 & a_{m-1,n-1}x^{m-1} y^{n-1} + \dots + a_{1,1}xy + a_{1,0}x + a_{0,1}y + a_{0,0}]. \tag{12}
 \end{aligned}$$

式(12)就是给出的一般二元切触有理插值公式,通过选取不同次数的多项式 $q(x,y)$ 可以得到不同次数类型的有理插值函数,相比其它算法具有下列优点:

- 1) 式(12)的分母多项式,可以是在矩形网格(1)上,函数值不为0的任意属于 $P_{2s+1,2t+1}$ 的二元代数多项式,这样就可以大大降低分母多项式的次数.
- 2) 其它方法计算分母多项式需要很大的计算量,而本文是直接给出的显式表达式.
- 3) 本文无论分母多项式次数是多少,分子多项式中的基函数 $\alpha_i(x), \beta_i(x), \alpha_j(y), \beta_j(y)$ 及 $p_{ij}(x,y)$ 始终不变,这样当改变有理函数的次数类型时,就减少了很大的计算量.

下面讨论向量值切触有理插值函数的算法.引入插值算子如下:

$$N_{ij}(x,y) = (x - x_i)(y - y_j) \mathbf{V}_{ij}^{(1,1)} + (x - x_i) \mathbf{V}_{ij}^{(1,0)} + (y - y_j) \mathbf{V}_{ij}^{(0,1)} + \mathbf{V}_{ij}. \tag{13}$$

显然

$$N_{ij}(x_i, y_j) = \mathbf{V}_{ij}, \quad \frac{\partial N_{ij}(x_i, y_j)}{\partial x} = \mathbf{V}_{ij}^{(1,0)}, \quad \frac{\partial N_{ij}(x_i, y_j)}{\partial y} = \mathbf{V}_{ij}^{(0,1)}, \quad \frac{\partial^2 N_{ij}(x_i, y_j)}{\partial x \partial y} = \mathbf{V}_{ij}^{(1,1)}.$$

定理 3 给定插值节点(1)和相应向量值 \mathbf{V}_{ij} 及向量偏导数值 $\mathbf{V}_{ij}^{(u,v)}$, 则

$$\mathbf{R}(x,y) = \sum_{i=0}^s \sum_{j=0}^t \lambda_{ij}(x,y) N_{ij}(x,y) \tag{14}$$

是满足插值条件的二元向量值切触有理插值函数.

事实上,将定理 1 中的函数换成向量,采用类似的方法即可证明.同理可得

$$\begin{aligned}
 \mathbf{R}(x,y) = & \sum_{i=0}^s \sum_{j=0}^t [q_{ij}\alpha_i(x)\alpha_j(y) + q_{ij}^{(1,0)}\alpha_j(y)\beta_i(x) + q_{ij}^{(0,1)}\alpha_i(x)\beta_j(y) + \\
 & q_{ij}^{(1,1)}\beta_i(x)\beta_j(y)] \mathbf{N}_{ij}(x,y) \Big/ [a_{m,n}x^m y^n + a_{m,n-1}x^m y^{n-1} + \\
 & a_{m-1,n}x^{m-1} y^n + a_{m-1,n-1}x^{m-1} y^{n-1} + \dots + a_{1,1}xy + a_{1,0}x + a_{0,1}y + a_{0,0}]. \tag{15}
 \end{aligned}$$

式(15)就是给出的一般二元向量值切触有理插值公式.通过选取不同次数的多项式 $q(x,y)$ 可以得到不同的次数类型的二元向量值有理插值函数,且相比其它算法同样具有上述的几个优点.

2 数值例子

通过下面的数值例子可以验证本文有理插值公式的算法可行性是无条件的,具有计算量较低、有理函数次数较低等特点.

例 1 已知

$$\begin{aligned} x_0 = 0, x_1 = 1, y_0 = 0, y_1 = 1, f_{00} = 0, f_{00}^{(0,1)} = 0, f_{00}^{(1,0)} = 1, f_{00}^{(1,1)} = 3, \\ f_{01} = 1, f_{01}^{(0,1)} = 1, f_{01}^{(1,0)} = 2, f_{01}^{(1,1)} = -5, f_{10} = -1, f_{10}^{(0,1)} = -1, f_{10}^{(1,0)} = 1, \\ f_{10}^{(1,1)} = -1, f_{11} = 0, f_{11}^{(0,1)} = 0, f_{11}^{(1,0)} = 1, f_{11}^{(1,1)} = 4, \end{aligned}$$

试用式(12)求分母多项式次数为(1,0)的二元切触有理插值函数,使之满足上述插值条件.

解 根据题意,由 Hermite 插值基函数可得

$$\begin{cases} \alpha_0(x) = 2x^3 - 3x^2 + 1, \alpha_1(x) = -2x^3 + 3x^2, \\ \beta_0(x) = x^3 - 2x^2 + x, \beta_1(x) = x^3 - x^2; \\ \alpha_0(y) = 2y^3 - 3y^2 + 1, \alpha_1(y) = -2y^3 + 3y^2, \\ \beta_0(y) = y^3 - 2y^2 + y, \beta_1(y) = y^3 - y^2. \end{cases}$$

由式(8)的插值算子公式得

$$\begin{cases} p_{00}(x, y) = 3xy + x, p_{01}(x, y) = -5xy + 7x + y, \\ p_{10}(x, y) = -xy + x - 2, p_{11}(x, y) = 4xy - 3x - 4y + 3, \end{cases}$$

根据式(12)得

$$\begin{aligned} q(x, y) &= x + 1, \\ r(x, y) &= \sum_{i=0}^1 \sum_{j=0}^1 [q_{ij} \alpha_i(x) \alpha_j(y) + q_{ij}^{(1,0)} \alpha_j(y) \beta_i(x) + \\ &\quad q_{ij}^{(0,1)} \alpha_i(x) \beta_j(y) + q_{ij}^{(1,1)} \beta_i(x) \beta_j(y)] p_{ij}(x, y) / q(x, y) = \\ &\quad [(78y^4 - 177y^3 + 90y^2 + 12y)x^4 + \\ &\quad (-160y^4 + 370y^3 - 195y^2 - 20y + 6)x^3 + \\ &\quad (66y^4 - 161y^3 + 93y^2 + 3y - 9)x^2 + \\ &\quad (14y^4 - 33y^3 + 18y^2 + 3y + 1)x + (-2y^4 + 3y^3)] / (x + 1). \end{aligned}$$

例 2 已知

$$\begin{aligned} x_0 = 0, x_1 = 1, y_0 = 0, y_1 = 1, \\ \mathbf{V}_{00} = (0, 1), \mathbf{V}_{00}^{(0,1)} = (0, -2), \mathbf{V}_{00}^{(1,0)} = (1, 5), \mathbf{V}_{00}^{(1,1)} = (-1, -4), \\ \mathbf{V}_{01} = (1, 3), \mathbf{V}_{01}^{(0,1)} = (1, -6), \mathbf{V}_{01}^{(1,0)} = (2, -3), \mathbf{V}_{01}^{(1,1)} = (0, 4), \\ \mathbf{V}_{10} = (-1, 2), \mathbf{V}_{10}^{(0,1)} = (-1, 4), \mathbf{V}_{10}^{(1,0)} = (1, -5), \mathbf{V}_{10}^{(1,1)} = (-1, 3), \\ \mathbf{V}_{11} = (0, 6), \mathbf{V}_{11}^{(0,1)} = (0, -2), \mathbf{V}_{11}^{(1,0)} = (1, 3), \mathbf{V}_{11}^{(1,1)} = (0, -5), \end{aligned}$$

试用式(15)求分母多项式次数为(1,0)的二元向量值切触有理插值函数,使之满足上述插值条件.

解 根据题意,由 Hermite 插值基函数可得

$$\begin{cases} \alpha_0(x) = 2x^3 - 3x^2 + 1, \alpha_1(x) = -2x^3 + 3x^2, \\ \beta_0(x) = x^3 - 2x^2 + x, \beta_1(x) = x^3 - x^2, \\ \alpha_0(y) = 2y^3 - 3y^2 + 1, \alpha_1(y) = -2y^3 + 3y^2, \\ \beta_0(y) = y^3 - 2y^2 + y, \beta_1(y) = y^3 - y^2. \end{cases}$$

由插值算子式(13)得

$$\begin{cases} N_{00}(x, y) = (-xy + x, -4xy + 5x - 2y + 1), \\ N_{01}(x, y) = (2x + y, 4xy - 7x - 6y + 9), \\ N_{10}(x, y) = (-xy + x - 2, 3xy + y - 5x + 7), \\ N_{11}(x, y) = (x - 1, -5xy + 8x + 3y - 2). \end{cases}$$

根据式(15)得

$$q(x, y) = x + 1,$$

$$\mathbf{R}(x, y) = (R_1(x, y), R_2(x, y)) =$$

$$\sum_{i=0}^1 \sum_{j=0}^1 [q_{ij} \alpha_i(x) \alpha_j(y) + q_{ij}^{(1,0)} \alpha_j(y) \beta_i(x) + q_{ij}^{(0,1)} \alpha_i(x) \beta_j(y) + q_{ij}^{(1,1)} \beta_i(x) \beta_j(y)] N_{ij}(x, y) / q(x, y),$$

$$R_1(x, y) = [(-6y^3 + y^2)x^4 + (-6y^4 + 25y^3 - 24y^2 + 6)x^3 + (8y^4 - 24y^3 + 18y^2 - y + 9)x^2 + (-4y^4 + 4y^3 + 3y^2 - y + 1)x + (-2y^4 + 3y^3)] / (x + 1),$$

$$R_2(x, y) = [(-96y^4 + 294y^3 - 225y^2 - 21y + 30)x^4 + (196y^4 + 646y^3 + 582y^2 + 26y - 68)x^3 + (-76y^4 + 308y^3 - 291y^2 + 11y + 35)x^2 + (-8y^4 + 20y^3 - 12y^2 - 6y + 6)x + (8y^4 - 28y^3 + 24y^2 - 2y + 5)] / (x + 1).$$

从上面例子可以看到,有理函数是满足插值条件的,分母多项式次数可以降低到任意次数(x 的次数为1, y 的次数为0).至于如何选取分母多项式的次数,则要视情况而定,特别是当插值节点的数量较大时,本文的方法更能降低分母的次数.

3 结 论

本文利用 Hermite 插值基函数和二元多项式插值误差的性质,给出的只是一种针对一阶导数情形的二元(向量值)切触有理插值算法.较之其它方法不仅具有较低计算量,而且算法的可行性也是无条件的,特别是当插值节点的数量较大时,本文算法的简便性更能够得到体现.但对于更高阶导数情形的二元(向量值)切触有理插值算法并没有解决,如何选取分母多项式,使得二元(向量值)切触有理插值函数的计算量最小、逼近效果最佳,并没有提及.这些问题值得读者进一步研究.

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Bivariate Osculatory Rational Interpolation on Rectangular Grids

JING Ke^{1,2}, KANG Ning²

(1. *School of Mathematics and Statistics, Fuyang Teachers College,*

Fuyang, Anhui 236037, P.R.China;

2. *School of Management, Hefei University of Technology,*

Hefei 230009, P.R.China)

Abstract: The bivariate osculatory rational interpolation is an important element of rational interpolation, and reducing the degrees of the osculatory rational interpolation functions and solving their existence make an important problem. The bivariate osculatory rational interpolation algorithms mostly have conditional feasibility and massive computational complexity with high function degrees. A bivariate osculatory rational interpolation algorithm was obtained on rectangular grids and extended to vector-valued cases, with the method of bivariate Hermite interpolation basis function in view of the error characteristics of bivariate polynomial interpolation. The numerical examples illustrate that, compared to other methods, the feasibility of the presented algorithm is unconditional, the degrees of the related rational functions are lower, and the algorithm has less computational complexity.

Key words: bivariate osculatory rational interpolation; error formula; bivariate Hermite interpolation; basis function

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